

Redes de Comunicação em Ambientes Industriais Aula 13

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In the last episode ... (1)

CANopen

- CAN High Layer Protocol
- Object-oriented Modeling of Device and Network
- Interoperability between Devices
- Interchangeability of Devices
- Off-the-shelf Plug-and-play Capability
- Off-the-shelf Configuration and Analysis Tools
- Standardized Communication Services
 - Peer-to-peer communication
 - Segmented Data Transfer
- Network and Node Configuration
- Network and Node Error Handling



In the last episode ... (2)

CANopen

Protocols:

- Process Data Object (PDO) Protocol
- Service Data Object (SDO) Protocols
- Synchronization (SYNC) Protocol
- Emergency (EMCY) Protocol
- Network Management Protocols:
 - □ NMT Message Protocol
 - Boot-Up Protocol
 - Error Control Protocol



standards.ieee.org/getieee802/802.3.html



- Created in the mid 70s (!) by Robert Metcalfe at the Xerox Palo Alto Research Center.
- Aimed initially at sharing expensive computing peripherals in office environment (particularly, it was developed to connect a printer to a computer)
- From then on, it <u>became extremely popular</u>, being used in <u>many</u> <u>domains beyond</u> its original scope. Particularly in <u>industrial</u> <u>systems</u>, it became the most serious candidate for the unifying communications protocol, from the plant floor to the management.
- ✓ Standardized in the mid 80s as IEEE 802.3

- Multi-master, broadcast, serial bus (initially) or star (currently)
 - ✓ <u>Bus</u>: ☺ simpler cabling / ⊗ faults in any point may cause complete communication disruption; fault tracing difficult
 - ✓ <u>Star</u>: ⊗ cabling more complex and expensive / ☺ more fault tolerant, easier and faster fault source tracing
- Synchronous transmission with Manchester bit coding
- Transmission rate of 10, 100Mbit/s, 1 and 10Gbit/s
- Max. number of nodes is 1024 (normally limited by a lower number of ports of the networking equipment)
- Max. link length is 100m (100Mbit/s, UTP cat5 cabling)
- Max. of 2 hubs between any 2 nodes (100Mbit/s)
- 2 architectures: shared (hubs), segmented (switches)
- Addressing classes: unicast, multicast and broadcast



Frame structure

Dst. Ac	dr	Src. Addr	Type/Len	LLC	SNAP	Data	CRC
6		6	2	3	5	<mark>38-1492</mark> 46-1500	4

LLC defined in IEEE802.2

LLC: Logic Link Control (Destination & Source Service Access Point + Control) SNAP: Sub-Network Access Protocol (used for IP packets)

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- Collisions can occur during the interval of one *slot* after start of \checkmark transmission (512 bits)
- When a collision is detected a **jamming** signal is sent (32 bits long) \checkmark
- Frames vary between 64 (min) and 1518 (max) octets physically add 7+1 octets for preamble & SOF and 96 bit times for IFS RCAI 2005/2006



- Back-off and Retry mechanism:
 BEB = Binary Exponential Backoff
- Retry instant in number of slots is randomly chosen (uniform distribution) within randomization window.
- Randomization window starts with [0, 2) slots and duplicates every retry until [0,1024) slots.
 Maximum of 16 retries.
- Performance degrades after ~60% utilization
 - For higher utilization it is prone to thrashing caused by chained collisions (not suited for real-time behavior)



Thrashing effect

As the offered load increases, the actual load handled by the network (throughput) decreases



Maximum network-induced delays can be very long with high packet-drop rates



Switched Ethernet

Ethernet switches



- Nowadays the most common solution (even for general purpose networks)
- No collisions; nodes see a private collision domain (microsegmented architecture)
- Messages directed to a busy output port are temporarily stored in the switch memory
- Up to 8 priorities



Switched Ethernet

Using switches

- Became the most common solution
 - Current switches are wire-speed (non-blocking)
 - 802.1D possibly with multiple priority queues (802.1p)
 - 802.1Q Virtual LANs

Vot perfect !

- Priority inversions in queues (normally FIFO)
- Mutual interference through shared memory and CPU
- Additional forwarding delay (with jitter caused by address table look up, address learning, flooding)
- Delays vary with switch technology and internal traffic handling algorithms



Why Real-Time Ethernet

 Motivations for using Ethernet in application domains beyond its original one e.g. factory automation, large embedded systems

(Decotignie, 2001)

- Cheap wrt other high speed technologies
- Widely available
- Scalable tx rate
- High bandwidth
- Easy integration with office networks (important for logging, management and multi-level integration, e.g. CIM)
- IP stacks widely available
- Well known and mature technology



Why not Real-Time Ethernet

 Downside of using Ethernet in application domains beyond its original one e.g. factory automation, large embedded systems

- Connection costs higher than traditional fieldbuses (PHY, transformer, MAC)
- High communication overhead (for short data items)
- High computing power requirements (for efficient bandwidth usage)
- Typical star topology not always adequate (may lead to extra long cabling)
- Existing protocol stacks (mainly IP) not optimized for real-time operation

⁽Decotignie, 2001)



Making Ethernet real-time

- ✓ Keep network load low (~1%, e.g. NDDS) !
- Avoid bursts (e.g. traffic smoothing)
- Modify the back-off and retry mechanism (e.g. CSMA/DCR, Virtual-time, windows, EQuB)
- Control transmission instants (e.g. token-passing (RETHER, RT-EP), master-slave (FTT-Ethernet, ETHERNET-Powerlink), TDMA (PROFINET),...)
- Micro-segmented Switched Ethernet with admission and transmission control



CSMA/CD based.

E.g. ORTE, Network Data Delivery Service (NDDS)

- Statistical guarantees based on limitation of the bandwidth utilization
- Publisher/subscriber cooperation model
- NDDS database: publishers and subscribers metadata
- NDDS database is shared among all nodes (holistic view)
- NDDS Library: set of functions accessible to the application to
 - System set-up, Update/get variable's values, etc.
- NDDS tasks
 - send and receive publication updates, NDDS database updates, etc.





CSMA/CD based.

E.g. Traffic-smoothing

- Statistical guarantees based on limitation of the traffic bandwidth and burstiness
- Real-time traffic transmitted immediately
- Non-real-time traffic load and burstiness controlled using a shaper (leacky bucket)
- Adaptive techniques recently proposed
 - Shaper parameters addapted according to the instantaneous communication characteristics (load and number of collisions)





Modified CSMA/CD protocols

- E.g. CSMA/DCR, Virtual-time, EQuB
 - Upon collisons RT nodes send black bursts (i.e. non-standard <jam signals) with length dependent on their priority (e.g. waiting time)
 - Bus state continuosly monitored. When a node senses no contention initiates the data transmission immediately





Token passing
 E.g. RT-EP, RETHER

- NRT NRT NRT NRT NRT NRT 1 2 2 3 3 4 4 5 5 6 6 RT 1 MRT 1 MRT
- Real-time data assumed as periodic
- Bus time divided in cycles of fixed duration
- Access to the bus regulated by the token
- First all RT nodes are visited
- Until the end of the cycle the token visits NRT sources
- ✓ E.g.
 - ✓ cycle *i* {1 − 4 − 1 − 2 − 3 − 4 − 5 − 6}
 - ✓ cycle *i*+1 {1 − 4 − 1 − 2},
 - ✓ cycle *i*+2 {1 − 4 − 1 − 2 − 3 − 4}
- Explicit RT requests, online admission control



Master/Slave
 E.g. ETHERNET
 Powerlink



- Fixed duration cycles
- Four phases (Start, Cyclic, Asynchronous and Idle)
- Periodic (isochronous) and event (asynchronous) data exchanges
- Network Manager (NM) (Master) coordinates the access to the bus
- Powerlink controllers (Slaves) transmit only after explicit request of the NM
- Different implementations allowed
 - Software
 - Dedicated CPU
 - Hardware
- Hub (high-speed) and switch architectures allowed

Time





- Two phases (Synchronous and Asynchronous windows)
- Periodic (isochronous) and event (asynchronous) data exchanges
- Master/Multislave . Master coordinates the access to the bus via a so-called trigger message, which contains the EC-schedule
- Slaves transmit only after explicit request of the Master
- Online scheduling
- Online admission control
- Arbitrary scheduling policies supported
- Hub and switch architectures allowed



- Switched Ethernet
 E.g. Ethernet/IP,
 (modified) EDF switch
- Real-time and non-real-time traffic





- RT traffic scheduled according to EDF and subject to admission control
- Based on RT layer (RT-I) both in nodes and (modified) Ethernet switches
- 🗸 RT-I
 - RT connection set-up
 - Admission control (both up and down links)
 - Deadline partitioning between the up and downlink
 - Message transmission and reception



Schedulability analysis

- Shared Ethernet: depends on the specific overlay protocol
 - Master-Slave, TDMA, Token-passing, Traffic smoothing
- Switched Ethernet: depends on queuing policies and switch internals
 - There are models that can be used directly, including for multiple priority queues and chained switches (e.g. the Network Calculus)
- For high bit rates and especially IP communication the delays in the protocol stacks become more important than the network-induced delays!



Schedulability analysis (switched Ethernet)

- Requires knowledge about the incoming traffic
- Typical analysis for fixed-priorities:
 - Network Calculus (Cruz, 1991)
 also known as the (σ,ρ) model
 - Response-time analysis

 (similar to CAN as previously shown)





Schedulability analysis (switched Ethernet – Network Calculus)

- ✓ Cumulative arrival at queue *i* (*F_i*) upper bounded by (σ_i, ρ_i): *F*(*t*) - *F*(*s*) ≤ $\sigma_i + \rho_i^*(t-s)$ $\forall_{0 \le s \le t}$
- Decreasing priorities with i
- ✓ Channel capacity c
- Upper bound on delay at queue i D_i





- Schedulability analysis
 (switched Ethernet Network Calculus)
 - One advantage of the network calculus is that is allows determining a bound to the burstiness of the outgoing traffic (σ'_i) and consequently to the **buffer requirements** of that flow

$$\sigma'_{i} = \sigma_{i} + \rho_{i} * \frac{\sum_{j=1}^{i-1} \sigma_{j} + \max_{i+1 \le j \le N} (C_{j})}{c - \sum_{j=1}^{i-1} \rho_{j}}$$

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Summary

- Ethernet frame structure (minimum data length is 46 octets)
- Bus and star topologies
- CSMA/CD operation
 - Non-deterministic collision-resolution mechanism
- Shared vs Switched Ethernet
- Why (not) Ethernet for Real-Time applications
- (Very!) Brief review of Ethernet-based RT protocols
 - CSMA/CD based: NDDS, Traffic Smoothing
 - Modified CSMA/CD: EQuB
 - Token-passing: Rether
 - ✓ Master/Slave: ETHERNET Powerlink, FTT-Ethernet (the best ☺)
 - Switched Ethernet: EDF switch
- Analitical tools
 - Network calculus